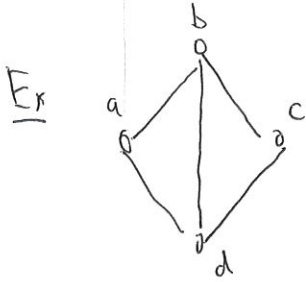


[Leskovec et al, Chapter 10.2, 10.4]

Betweenness Centrality : Measurement how much importance a link is

- All pairs of nodes have 1 communication.
- They will use a shortest path for a communication.
- If there are more than one shortest path, they will choose the communication path randomly.



communication between  $a \rightarrow c$

2 shortest paths

$a \rightarrow b \rightarrow c \rightarrow$  use with prob ~~1~~ 0.5

$a \rightarrow d \rightarrow c \rightarrow$  use with prob 0.5

- Count expected number of communications on each link.

Ex expected #communications on link  $\{a,b\}$

- from $a \rightarrow c$	$\rightarrow$ prob. 0.5
- from $a \rightarrow b$	$\rightarrow$ prob. 1
- from $a \rightarrow d$	$\rightarrow$ prob. 0
- from $b \rightarrow a$	$\rightarrow$ prob. 1
- from $b \rightarrow c$	$\rightarrow$ prob. 0
- from $b \rightarrow d$	$\rightarrow$ prob. 0
- from $c \rightarrow a$	$\rightarrow$ prob. 0.5
- from $c \rightarrow b$	$\rightarrow$ prob. 0
- from $c \rightarrow d$	$\rightarrow$ prob. 0
- from $d \rightarrow a$	$\rightarrow$ prob. 0
- from $d \rightarrow b$	$\rightarrow$ prob. 0
- from $d \rightarrow c$	$\rightarrow$ prob. 0

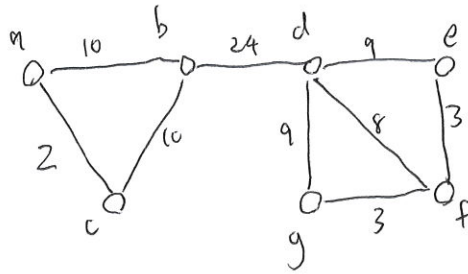
Expected #communications  
 = betweenness of  $\{a,b\}$   
 = 3

Betweenness of  $\{a, d\}, \{b, c\}, \{c, d\}$  is 3.

Betweenness of  $\{b, d\}$  is 2.

Formulation: Betweenness of  $\{u, v\} = \sum_{x, y \in V} \frac{\# \text{ shortest paths using } \{u, v\} \text{ from } v_1 \rightarrow v_2}{\# \text{ shortest paths from } v_1 \rightarrow v_2}$

Exercise



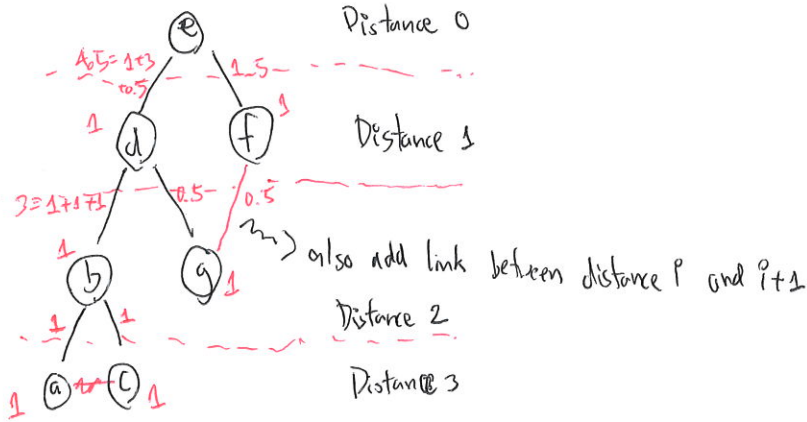
Algorithms for calculating betweenness centrality

Technique: Breadth First Search (BFS)

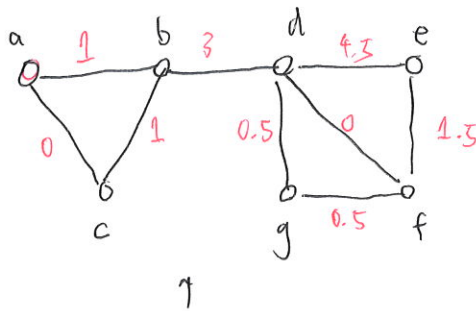
Use for calculating # shortest paths from  $e$  to all other nodes.

Question

How to calculate # shortest paths from  $e$  to other nodes?



BFS Tree



- Calculate the graph like this, but begin from all other nodes.
- Add the results together.

# Algorithm

For all node  $i$ :  $O(|V||E|)$

Perform a BFS to construct a BFS tree from  $i$

$O(|E|)$

Consider the BFS Bottom-up, Suppose that we are considering  $j$  at level  $d$ .

$$cen_{i,j}^d = \sum_{j': \text{incident edge at level } d+1} [cen_{i,j'}^{d+1}] + 1$$

# incident edges at level  $d$ .

Same for the edge incident to the same node

$O(|E|)$

Can be pre-calculated at nodes!

Betweenness of edge  $j = \sum_p cen_{i,j}^p$

Computation Time  $O(|V||E|)$

## Graph Clustering Method (Divide a social network into groups)

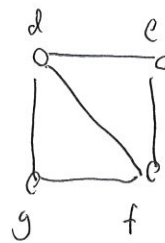
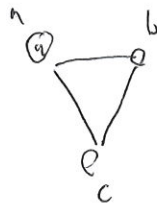
1: Remove the edge with highest betweenness centrality

2: If "nice groups" are obtained, terminated.

Otherwise, go to step 1

3: Nodes that can communicate together (after the removals) will be put into the same group.

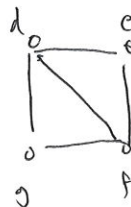
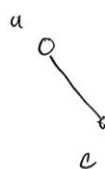
Remove betweenness  $\geq 12$



Group 1: {a, b, c}

Group 2: {d, e, f, g}

Remove betweenness  $\geq 5$



Group 4: {a, c}

Group 2: {b}

Group 3: {d, e, f, g}

Create groups by removing as least as edges as possible

# Partition a social network [k-way clustering]

o Divide a social network into 2 groups with similar sizes,

## Laplacian Matrix

$$A = \begin{matrix} & \begin{matrix} a & b & c & d & e & f & g \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \\ f \\ g \end{matrix} & \begin{pmatrix} 2 & -1 & -1 & 0 & 0 & 0 & 0 \\ -1 & 3 & -1 & -1 & 0 & 0 & 0 \\ -1 & -1 & 2 & 0 & 0 & 0 & 0 \\ 0 & -1 & 0 & 4 & -1 & -1 & -1 \\ 0 & 0 & 0 & -1 & 2 & -1 & 0 \\ 0 & 0 & 0 & -1 & -1 & 3 & -1 \\ 0 & 0 & 0 & -1 & 0 & -1 & 2 \end{pmatrix} \end{matrix}$$

$A = (a_{ij})$  where

$$a_{ij} = \begin{cases} 0 & \text{when } i \neq j \text{ and } \{i, j\} \notin E \\ -1 & \text{when } i \neq j \text{ and } \{i, j\} \in E \\ \text{degree of } i & \text{when } i = j \end{cases}$$

Find vector  $v$  and value  $\lambda$  such that  $\lambda v = A \cdot v$

Eigenvalue decomposition

Smallest eigenvalue = 0

eigenvector =  $\begin{bmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix}$

$$0 \cdot \begin{bmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \begin{bmatrix} a \\ \vdots \\ g \end{bmatrix} = \begin{bmatrix} a & \dots & g \\ 2 & -1 & -1 & 0 & 0 & 0 \\ -1 & 3 & -1 & -1 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ g & \dots & \dots & \dots & \dots & \dots \end{bmatrix} \cdot \begin{bmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix}$$

$A \quad \cdot \quad v$

This is true for all Laplacian matrix

Second smallest eigenvalue = 0.398

eigenvalue =  $\begin{bmatrix} a \\ b \\ c \\ \hline d \\ e \\ f \\ g \end{bmatrix} = \begin{bmatrix} -0.493 \\ -0.297 \\ -0.493 \\ \hline 0.214 \\ 0.356 \\ 0.356 \\ 0.356 \end{bmatrix}$

→ one partition

→ another partition

